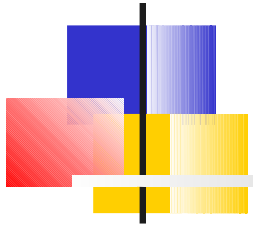


Performance Architecture exercises

P. Bakowski



bako@ieee.org





Ex1: Basic performance

- The **execution time** for a given task is a reliable method to evaluate computer performance.

The execution time is given by the following equation:

$$\text{ex_time} = \frac{\text{instructions/task} \times \text{clock_cycles/instruction}}{\text{seconds/clock_cycle}}$$



Ex1: Basic performance

1. **clock_cycle/instruction** must be considered as an **average value (explain why ?)**

2. Let us calculate **ex_time** for:

instructions/program=**100 Gi**

clock_cycles/instruction = **1.2**

clock frequency = **2.4 GHz**



M-Ex1: Basic performance

instructions/program=100 Gi,

clock_cycles/instruction = 1.2 and

clock frequency = 2.4 GHz

$$\text{ex_time} = 100 * 10^9 * 1.2 * 1/2.4 * 10^{-9} =$$

$$100 * 1.2 / 2.4 = 50 \text{ s}$$



Ex2: Performance in MIPS

- Let us consider n instruction classes [max 4] and the corresponding **CPI** measurements (Clock Cycles per Instruction) [from 1 to 4 CPI] for two computers.
- Each measurement is relative to the same program.
- Assume that the computer clock rate is $CR = 1$ GHz.
- Which code sequence will execute faster according to MIPS and according to execution time ?



Ex2: Performance in MIPS

Number of CPI classes per computer, in each case we give the number of G_i in the class with $CR = 1$ GHz:

$$NG_{iA} = [1, 2, 5, 4]$$

$$NG_{iB} = [3, 6, 1, 3]$$

For each class the **CPI** is given as:

$$CPI = [1, 2, 3, 4]$$

- What is the execution time for each program (A,B) ?
- What is the MIPS measure for A and B ?



M-Ex2: Performance in MIPS

$NGiA = [1, 2, 5, 4]$, $NGiB = [3, 6, 1, 3]$

$CPI = [1, 2, 3, 4]$,

$CR = 1 \text{ GHz}$

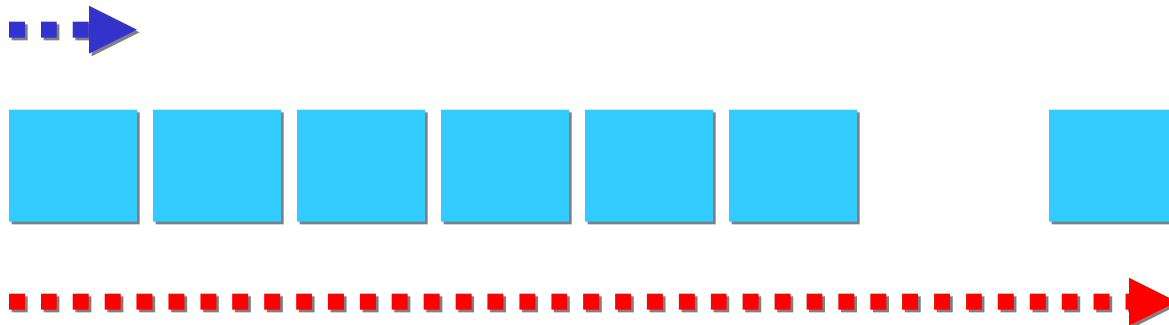
$extA = 1*1*/CR + 2*2/CR + 5*3/CR + 4*4/CR = 36/CR$

$extB = 3*1*/CR + 6*2/CR + 1*3/CR + 3*4/CR = 30/CR$

- What is the execution time for each program (A,B) ?
- What is the MIPS measure for A and B ?

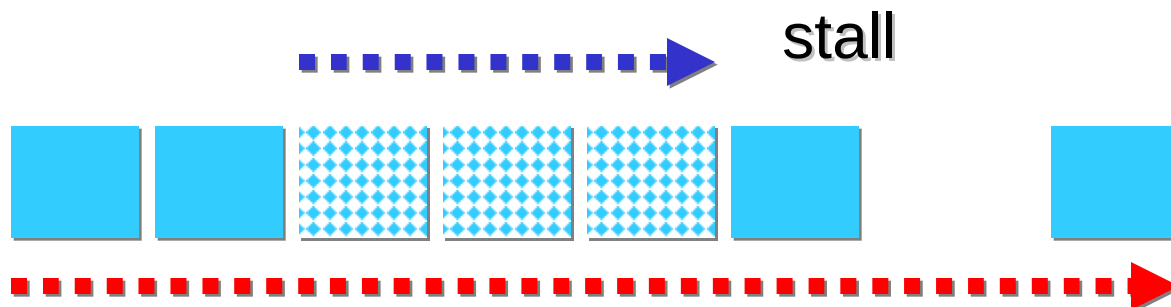
Ex3: Pipeline performance

- In this exercise we evaluate the **execution speed** up for a given pipeline.
- Each pipeline stage correspond to one clock cycle.
- The length of the pipeline is a parameter.



Ex3: Pipeline performance

- The executed program is a mix of instructions:
 - calculation instructions may be executed without pipeline stall
 - **conditional branch** instructions involve **bps** stages of delay (stall)
 - **memory access** instruction involve **mps** stages of delay





Ex3: Pipeline performance

- Given the clock frequency, the number of stages, the number of instructions to execute and the **proportion** of **conditional**/**memory_access**/**calculation** instructions,

we can evaluate the **necessary execution time** for non-pipelined and pipelined architectures.



Ex3: Pipeline performance

parameters

- clock frequency in MHz: 1000
- number of pipeline stages: 8
- total instruction number in millions (inb): 1000
- ratio of branch instructions : 0.2
- average stall **penalty** for branch instructions : 6
- ratio of memory access instructions : 0.2
- average stall **penalty** for memory access instructions : 4



Ex3: Pipeline performance

Non pipelined execution time is given as:

$$\text{npipe_time} = \text{clock_cycle} * \text{number_of_stages} * \text{number_of_instructions}$$

Pipelined execution time is calculated as:

$$\text{calrt} = 1.0 - \text{brst} - \text{memrt}$$

caltr - ratio of calculation instructions

$$\text{pipe_time} = \text{clock_cycle} * \text{inb} * \text{caltr} + \text{clock_cycle} * \text{inb} * \text{brstall} * \text{brst} + \text{clock_cycle} * \text{inb} * \text{memstall} * \text{memrt}$$



M-Ex3: Pipeline performance

Non pipelined execution time is given as:

$$\text{npipe_time} = 10^{-9} * 8 * 10^9 = 8 \text{ s}$$

Pipelined execution time is calculated as:

$$\text{calrt} = 1.0 - \text{brrt} - \text{memrt} = 0.6$$

caltr - ratio of calculation instructions

$$\text{pipe_time} = 10^{-9} * 10^9 * 0.6 + 10^{-9} * 10^9 * 6 * 0.2 + 10^{-9} * 10^9 * 4 * 0.2 = 0.6 + 1.2 + 0.8 = 2.6 \text{ s}$$

Ex4: Power consumption

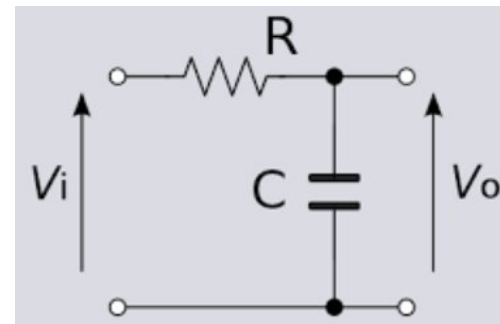
As we have seen in the performance equation, the easiest way to augment the performance is **to increase the frequency**.

The problem is here:

$$\text{dynamic.power} = A * N * C * V^2 * f$$

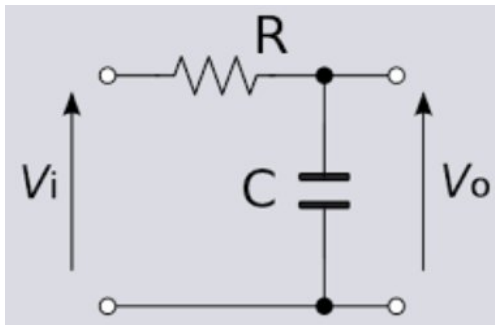
Apparently the dynamic power consumption (CMOS) is proportional to the frequency ?

But look at this ?



Ex4: Power consumption

$$\text{dynamic.power} = A * N * C * V^2 * f$$



For the given circuit, if we wish to get the same V_o twice as quickly we must increment V_i by the factor $x=?$

$$V_o = x * V_i * (1 - e^{-t/\tau})$$

Let us fix the ratio V_o/V_i to 0.8, and τ to 0.1ns. What is the maximum switching frequency ?

Now let us try, for the same circuit, to double the switching frequency, what is the x factor.



M-Ex4: Power consumption

Let us fix the ratio V_o/V_i to 0.8, and τ to 1ns. What is the maximum switching frequency ?

$$V_o = x * V_i * (1 - e^{-t/\tau}) \Rightarrow$$

$$0.8 = (1 - e^{-t/0.1}) \Rightarrow 0.2 = e^{-t/1}$$

$$\Rightarrow \ln 0.2 = -t/1 \Rightarrow$$

$$t = -1 * \ln 0.2 = 1.61 \text{ ns}$$

$$\text{frequency is } 1/1.61 \text{ ns} = 0.62 \text{ GHz}$$



M-Ex4: Power consumption

Now we wish to increase the frequency by the factor 2. So we need to reduce the t by the same proportion, it is $t=1.61$ ns \Rightarrow **0.8 ns**

$$V_o = x \cdot V_i \cdot (1 - e^{-t/\tau}) \Rightarrow$$
$$x = V_o / V_i \cdot (1 - e^{-t/\tau}) = 0.8 / (1 - e^{-t/\tau}) =$$
$$0.8 / (1 - e^{-0.8}) = 0.8 / (1 - 0.45) =$$
$$= 0.8 / 0.55 = 1.45$$

The resulting voltage increase x is: **1.45**



M-Ex4: Power consumption

Coming back to our initial equation:

$$\text{dynamic.power} = A * N * C * V^2 * f$$

Now we know that the increase of frequency by factor 2 involves the increase of the **voltage** by the factor – **1.45**

So the total power increase (for the same circuit) is:

$$1.45 * 1.45 * 2 = 2.12 * 2 = 4.24 \text{ times}$$



Ex5: Amdahl's law

The execution time of a mixed, sequential-parallel program, after the acceleration of the parallel part, is given by a simple equation known as **Amdahl's law**.

$$\text{tex_new} = \text{tex_par}/\text{acceleration} + \text{tex_seq}$$

Calculate the necessary acceleration to obtain the required new execution time.

Take:

$$\text{tex_par} = 10 \text{ s}, \text{tex_seq} = 5, \text{tex_new} = 8 \text{ s}$$



Ex5: Amdahl's law

$$\text{tex_new} = \text{tex_par}/\text{acceleration} + \text{tex_seq}$$

Calculate the necessary acceleration to obtain the required new execution time.

$$\text{tex_par} = 10 \text{ s}, \text{tex_seq} = 5, \text{tex_new} = 8 \text{ s}$$

$$\text{tex_par}/\text{acceleration} = \text{tex_new} - \text{tex_seq}$$

$$\text{acceleration} = \text{tex_par} / (\text{tex_new} - \text{tex_seq})$$

$$\text{acceleration} = 10 / (8 - 5) = 10 / 3 = 3.3$$



Summary

- Basic performance
- MIPS as performance measure
- Pipeline performance
- Power consumption
- Amdahl's Law